CaDiCaL(\mathcal{T}): CaDiCaL as CDCL(\mathcal{T}) Engine in cvc5

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- » state-of-the-art SMT solver
- $\,\gg\,$ most recent incarnation of the CVC tools
- » joint project led by Stanford University and University of Iowa in collaboration with
 - Universidade Federal de Minas Gerais (Brazil)
 - Bar Ilan University (Israel)
- » based on $CDCL(\mathcal{T})$ framework
- » supports wide range of theories in combination with quantifiers
 - all SMT-LIB theories + non-standard theories and theory extensions
- » support for **proofs** (incl. preprocessing, rewriting)
- » capabilities beyond standard SMT
 - SyGuS, abduction, interpolation, quantifier elimination, optimization (WIP)

- » propositional abstraction of the input formula
- $\hspace{0.1 cm} \hspace{0.1 cm} \hspace{0}$
- $\,\gg\,$ theory layer guides the search of the SAT solver
- » online, tight integration of SAT solver
 - theory layer interacts with SAT solver during the search
 - backward communication channel to notify theory layer about variable assignments, decisions, backtracks
 - theory layer derives conflicts, propagates theory literals, suggests decisions based on theory-guided heuristics

CDCL(T) SAT Solver: Current State-of-the-Art

- » no standardized SAT solver interface for interactive incremental SAT solving
- » solver-specific workarounds and modifications to the SAT solver
- » error prone, high potential for unintentional performance hits
- » difficult to replace
- » missed opportunities to take advantage of improvements in SAT

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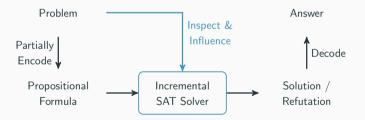
Situation in cvc5 (until recently)

- » integrates highly customized version of MiniSat
 - produces resolution proofs
 - push/pop for adding/deleting clauses and variables
 - custom theory-guided decision heuristics

IPASIR-UP in a Nutshell

IPASIR-UP = **IPASIR** + **U**ser **P**ropagators (Fazekas et al. 2023)

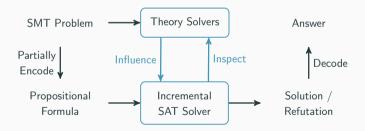
- » presented at SAT 2023
- » a SAT solver interface for
- » interactive incremental SAT solving



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- » presented at SAT 2023
- » a SAT solver interface for
- » interactive incremental SAT solving



» our focus: Integration as $CDCL(\mathcal{T})$ SAT solver

CaDiCaL(T) Integration (via IPASIR-UP)

- » Full utilization of interface
- » ~500 LOC in C++ for implementing interface (~800 including comments)
- » "easily" replaced with any SAT solver implementing IPASIR-UP
- » supports all* cvc5 features
- > *proof support work-in-progress
- » changes compared to MiniSat
 - resolution proofs \rightarrow DRAT/LRAT proofs (WIP)
 - \blacksquare native push/pop \rightarrow activation literals

IPASIR-UP Interface

Notifications (for Inspecting CDCL Search)

- > notify_assignment
- > notify_new_decision_level
- > notify_backtrack

Callbacks (for Influencing CDCL Search)

- ≫ cb_propagate
- > cb_add_reason_clause_lit
- > cb_decide
- > cb_add_external_clause_lit
- > cb_check_found_model

notify_assignment

- » sends assignment notification for observed variables
- » track assignment for theory literals
 - constructs (partial) assignment of propositional abstraction
- \gg track whether assignment is
 - decision
 - fixed
- » **notify** theory solvers about assigned theory literal,
 - e.g., if (observed) variable corresponding to theory literal a < 42 is assigned to
 - **true**: send a < 42 to arithmetic solver
 - **false**: send $a \ge 42$ to arithmetic solver

- » Used to manage the incremental state of cvc5
- » backtrackable data structures (context-dependent), associated with a context
 - SAT context, backtracks when SAT solver backtracks (decision-level push/pop)
 user context (SMT-level push/pop)

notify_new_decision_level

- » push SAT context
- » track decision + level

notify_backtrack (L)

- » pop SAT context to L
- **»** undo assignments at level > L
- » resend "popped" fixed theory literals
 - theory literals fixed at levels > L are popped
 - fixed assignments only notified once
 - resend fixed theory literals at level L

$cb_check_found_model$

- ≫ called when SAT solver found satisfying assignment
 ▷ returns true if assignment *T*-consistent, false otherwise
- » checks if assignment is *T*-consistent (full effort check)
 - \blacksquare theory solvers check $\mathcal T\text{-consistency}$ of assigned theory literals
 - » send conflict clause
 - » send lemmas
 - » send theory propagations
 - \triangleright adds eager explanations at this point to force SAT solver to propagate
 - *T*-consistent if:
 - » theory solvers performed all checks and
 - » no new variables were added and
 - » no new lemmas or conflicts sent, i.e., no new clauses added

$\textbf{cb_decide}$

- » called before SAT solver makes decision
- » used to inject theory-guided decisions
 - theory decisions (required)
 - $\,\vartriangleright\,$ decision strategies used by theory solvers
 - decision requests (optional)
 - \triangleright custom decision heuristics
 - \vartriangleright e.g.: justification heuristic, chooses next decision based on structure of formula
- » may discover partial satisfying assignment
 - triggers full effort check, i.e., calls cb_check_found_model
 - \blacksquare stops search if current assignment is $\mathcal{T}\text{-consistent}$

cb_add_external_clause_lit

- » clauses added during search are buffered
 - theory lemmas
 - theory conflicts
- » buffered clauses are only added during callback

cb_has_external_clause

» checks whether new clauses are pending

cb_propagate

- » called after SAT solver is done with propagation
- » performs lightweight checks in theory solvers (standard effort check)
- » theory propagations

cb_add_reason_clause_lit (prop_lit)

- » called when theory propagation prop_lit is involved in conflict
- » explain theory propagation
- » adds explanation (reason clause)

SMT push/pop via Activation Literals

- » happens between SAT solver calls, not during search
- » **push** assertion level
 - create fresh activation literal I_n for pushed level n
 - add I_n to each clause added in level n
 - prior to solving, assume $\neg l_i$ for $i \in \{1, n\}$
- » pop assertion level
 - add unit clause {*l_n*} for popped level *n* ▷ garbage collects all clauses added at level *n*
 - unobserve and fix value of variables introduced in n (important for performance)
- » renotify fixed literals with fixed level > intro level
 - requires keeping track of assertion levels when
 - » variable was introduced
 - » variable assignment was fixed

```
(set-logic ...)
```

. . .

. . .

(assert ...) ; A1 (assert ...) ; A2 (push 1) (assert ...) ; A3 (check-sat) (pop 1)

```
(check-sat)
```

» all incremental and non-incremental benchmarks of SMT-LIB 2023

- 434, 212 non-incremental benchmarks
- 43,287 incremental benchmarks
- » 300s time limit, 8GB memory limit
- » comparison of cvc5-1.0.8-dev with
 - MiniSat (custom, based on 2.2.0)
 - CaDiCaL (IPASIR-UP, version 1.7.4)

Evaluation: SMT-COMP Non-Incremental Divisions

	CVC5+MINISAT		CVC5+	CVC5+CADICAL	
Division	solved	time [s]	solved	time [s]	
Arith (6,925)	6,341	181,329	6,332	183,417	
BitVec (6,185)	5,645	168,844	5,625	175,110	
Equality (12,159)	5,331	2,060,608	5,337	2,059,279	
Equality+LinearArith (56,562)	45,970	3,196,706	45,966	3,198,129	
Equality+MachineArith (10,911)	1,073	2,958,372	1,075	2,958,746	
Equality+NonLinearArith (21,162)	13,333	2,425,551	13,123	2,474,917	
FPArith (3,979)	3,133	275,579	3,138	272,751	
QF_Bitvec (46,191)	43,735	1,092,892	43,713	1,092,907	
QF_Datatypes (8,403)	8,083	109,941	8,158	84,593	
QF_Equality (8,054)	8,043	9,338	8,047	6,968	
QF_Equality+Bitvec (16,801)	15,922	355,232	16,132	263,786	
QF_Equality+LinearArith (3,644)	3,464	65,242	3,497	52,176	
QF_Equality+NonLinearArith (906)	721	61,692	711	64,217	
QF_FPArith (76,252)	76,072	93,150	76,087	77,682	
QF_LinearIntArith (16,389)	11,530	1,604,847	12,017	1,489,186	
QF_LinearRealArith (2,008)	1,686	142,921	1,784	107,522	
QF_NonLinearIntArith (25,446)	13,076	4,080,649	14,058	3,696,580	
QF_NonLinearRealArith (12,134)	11,155	336,630	11,247	309,251	
QF_Strings (100,101)	98,407	619,928	98,870	483,260	
Total (434,212)	372,720	19,839,459	374,917	19,050,487	

- » +2197 solved instances
- $$\label{eq:solved_solved} \begin{split} & \sim 25\% \mbox{ faster on commonly} \\ & \mbox{ solved instances} \end{split}$$
- ≫ $2-4 \times$ faster in several logics
- » 13 of 19 divisions improved
 - quantifier-free better overall
 - quantified logics a bit behind
- » promising performance without much tuning or optimizations
- » solid baseline for future tuning with IPASIR-UP interface

	CVC5+MINISAT		CVC5+CADICAL	
Division	solved	time [s]	solved	time [s]
Arith (11)	41,362	233	41,362	240
BitVec (18)	36,114	2,992	36,117	3,031
Equality (4,067)	46,256	620,984	46,216	623,400
Equality+LinearArith (1,894)	431,172	57,390	430,552	59,637
Equality+MachineArith (4)	818	310	818	309
Equality+NonLinearArith (4,374)	82,721	651,804	83,801	644,742
FPArith (10)	3,422	1,849	3,421	1,849
QF_Bitvec (2,590)	51,334	63,165	51,260	62,036
QF_Equality (1,778)	29,981	4,616	29,982	4,588
QF_Equality+Bitvec (3,633)	7,677	148,084	7,620	153,446
QF_Equality+Bitvec+Arith (664)	959	51,776	985	44,466
QF_Equality+LinearArith (3,947)	2,266,894	130,331	1,893,335	133,167
QF_Equality+NonLinearArith (1,018)	96,917	24,307	92,813	23,932
QF_FPArith (19,188)	538,936	955,264	560,379	745,166
QF_LinearIntArith (69)	1,332,173	17,582	1,089,226	17,109
QF_LinearRealArith (10)	482	3,004	571	2,918
QF_NonLinearIntArith (12)	349,862	3,603	326,463	3,603
Total (43,287)	5,317,080	2,737,301	4,694,921	2,523,646

- » improvements in **some logics**
- » overall performance not there yet
- » poor performance on benchmarks with many check-sat calls
- » overhead of activation literals?

Observation

Performance poor on benchmarks with large number of check-sat calls

Example: kundu_true-*.smt2 (QF_LIA)

- 900k+ check-sat calls
- \gg solved queries within 300 seconds
 - MiniSat: 148,997
 - CaDiCaL: 103,843

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Activation Literal Overhead Experiment

- (push 1)...fresh literal I_n (assert true)...add clause $(I_n \vee \top)$ (check-sat)...assume $\neg I_n$ (pop 1)...add clause (I_n)
- > Repeated N times in one benchmark

Observation

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(push 1)	 fresh literal In
(assert true)	 add clause $(I_n \lor \top)$
(check-sat)	 assume $\neg I_n$
(pop 1)	 add clause (I_n)

» Repeated N times in one benchmark

N	MiniSat	CaDiCaL	Slowdown
10k	265ms	462ms	1.7 imes
25k	625ms	1.8s	2.8 imes
50k	1.2s	5.8s	4.8 imes
75k	1.8s	11.9s	6.6 imes
100k	2.5s	20.3s	8.1 imes

Conclusion

Summary

- » non-incremental performance solid
- » incremental performance still lagging behind
- » IPASIR-UP integration
 - simple and flexible
 - captures all functionality required by cvc5

What's Next?

- » DRAT/LRAT proof integration (WIP)
- » SAT solver tuning (currently default options)
- » improve performance on quantified problems
- » improve incremental performance
- » IPASIR-UP: reduce callbacks, notifications



https://cvc5.github.io

Scatter Plots QF_S*, QF_A*

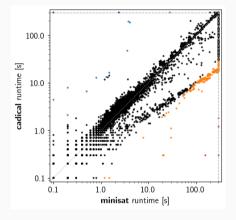


Figure 1: QF_S* (Logics with Strings)

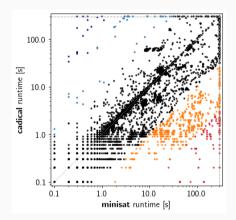


Figure 2: QF_A* (Logics with Arrays)

Quanitfier-free and Quanitifed Logics

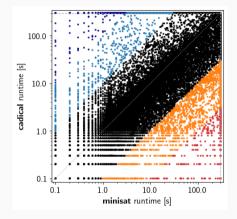


Figure 3: Quantifier-free Logics

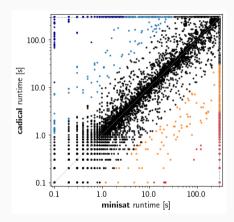


Figure 4: Quantified Logics

References

Fazekas, Katalin et al. (2023). "IPASIR-UP: User Propagators for CDCL". In: 26th International Conference on Theory and Applications of Satisfiability Testing, SAT 2023, July 4-8, 2023, Alghero, Italy. Ed. by Meena Mahajan and Friedrich Slivovsky. Vol. 271. LIPIcs. Schloss Dagstuhl - Leibniz-Zentrum für Informatik, 8:1–8:13. DOI: 10.4230/LIPIcs.SAT.2023.8. URL: https://doi.org/10.4230/LIPIcs.SAT.2023.8.