

An Invitation to Modal Logic: Lecture 4

Philosophy 150

Eric Pacuit

Stanford University
`ai.stanford.edu/~epacuit`

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Plan

- ✓ Motivating Examples
- ✓ Formalizing the muddy children puzzle, Introduction to Modal Logic
- ✓ More about truth of modal formulas

12/3: Focus on Epistemic Logic.

Digression: A small experiment.

12/5: Multiagent Epistemic Logic, Dynamics in Logic

12/7: Dynamics in Logic II

Plan for Today

First half of the lecture: Epsitemic Logic.

Second half of the lecture: a small experiment.

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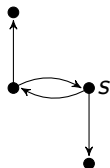
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Truth:

- ▶ $w \models \Box P$ iff for all v , if wRv then $v \models P$
- ▶ $w \models \Diamond P$ iff there exists v such that wRv and $v \models P$.

Two issues to remember

1. Modal formulas are interpreted **locally**.



$\mathbb{K}$



$\mathbb{N}$

No modal formula can distinguish between s and u .

Can you think of a first-order formula that can distinguish the Kripke structures?

Two issues to remember

1. Modal formulas are interpreted **locally**.
2. Modal logic can express interesting properties of Kripke structures.
 - $\Box P \rightarrow P$ **corresponds** to the **reflexivity** property.
 - $\Box P \rightarrow \Box \Box P$ **corresponds** to the **transitivity** property.

Next lecture: Dynamics in logic.

Questions?

Email: epacuit@stanford.edu

Website: ai.stanford.edu/~epacuit

Office: Gates 258